Deterministic and Stochastic, Time and Space Signal Models: An Algebraic Approach

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Structure and Digital Signal Processing

Is DSP algebraic?

- By restricting to Linear Algebra are we missing something?
- Apparently disparate concepts => instantiations same concept

Is DSP geometric?

- Constraints may restrict signals to a manifold
- Algorithms and signal processing should be derived for manifolds

Proposed Special Session for ICASSP'06

DSP: Algebra vs. Geometry

References for talk:

Pueschel and Moura, SIAM Journal of Computing, 35:(5), 1280-1316, March 2003 Pueschel and Moura, "Algebraic Theory of Signal Processing, 150 pages, Dec 2004



Algebraic Theory of SP

- Quick refresh on DSP
- DSP: Algebraic view point
 - Signal Model
- Algebraic Theory: Time
 - Time shift
 - Boundary conditions (finite time)
 - Fourier transforms, spectrum

Algebraic Theory: Space

- Space shift
- Infinite space: C-transform and DSFT
- Finite space: DTTs

What is it useful for:

- Fast algorithms
- m-D: separable and non-separable, new transforms



DSP

- Scalar, discrete index (time or space) linear signal processing
- 1-D or m-D: indexing set
 - Example: infinite discrete time
 - Signals:

$$S(z) = \sum_{n \in \mathbb{Z}} s_n z^{-n} \in \{\sum_{n \in \mathbb{Z}} s_n z^{-n} | (\dots, s_{-1}, s_0, s_1, \dots) \in \ell^2(\mathbb{Z})\} = \mathcal{M}$$

• Filters:

$$H(z) = \sum_{n \in \mathbb{Z}} h_n z^{-n} \in \{\sum_{n \in \mathbb{Z}} h_n z^{-n} | (\dots, h_{-1}, h_0, h_1, \dots) \in \ell^1(\mathbb{Z})\} = \mathcal{A}$$

Convolution (multiplication):

$$Y(z) = H(z) \cdot S(z) = \sum_{n \in \mathbb{Z}} \left(\sum_{i \in \mathbb{Z}} h_i s_{n-i} \right) z^{-n} \in \mathcal{M}$$

z-Transform:

$$\Phi: \qquad V = \ell^2(\mathbb{Z}) \to \mathcal{M}$$

$$\mathbf{s} = (\ldots, s_{-1}, s_0, s_1, \ldots) \quad \mapsto \quad \sum_{n \in \mathbb{Z}} s_n z^{-n}$$



DSP

Fourier Transform: DTFT

$$S(\omega) = S(e^{j\omega}) = \sum_{n \in \mathbb{Z}} s_n e^{-j\omega n}; \quad s_n = \frac{1}{2\pi} \int_{-\pi}^{\pi} S(\omega) e^{j\omega n} d\omega$$

- Spectrum: $(S(\omega))_{\omega \in W} = [-\pi,\pi)$
- Impulses: $E_{\omega}(z) = \sum_{n \in \mathbb{Z}} e^{j\omega n} z^{-n}$
- Eigen property: $H(z)E_{\omega}(z) = H(e^{j\omega n})E_{\omega}(z)$
- Linear combination: $\alpha H(z) + \beta H'(z)$; $\alpha S(z) + \beta S'(z)$



 ${\mathcal A}$ and ${\mathcal M}$ are vector spaces

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DSP: Algebraic View Point

• Cascading of filters: $H(z) \cdot H'(z)$

 \longrightarrow makes \mathcal{A} an algebra – the algebra of filters

• Convolution (multiplication): $Y(z) = H(z) \cdot S(z)$

Signal Model: Triplet $(\mathcal{A}, \mathcal{M}, \Phi)$

where bijective linear mapping $\,\, \Phi : \, V o \mathcal{M} \,$



DSP: Finite Time

- Signals: $S(z) = \sum_{i=0}^{n-1} s_i z^{-i} \in \mathcal{M} \stackrel{?}{=} \{\sum_{i=0}^{n-1} s_i z^{-i}\}$
- Filters: $H(z) = \sum_{i=0}^{n-1} h_i z^{-i} \in \mathcal{A} \stackrel{?}{=} \{\sum_{i=0}^{n-1} h_i z^{-i}\}$
- Convolution (multiplication): $Y(z) = H(z) \cdot S(z) \notin \mathcal{M}$

• Candidates: algebras of filters \mathcal{A} and modules of signals \mathcal{M} ?



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Algebraic Theory: Shift

Shift: special type of filter $x \in \mathcal{A}$

• $h \in \mathcal{A}$ expressed as polynomial or series in x

 $\blacksquare \implies x \text{ generates } \mathcal{A}$

Shift invariance: x shift, $\forall s, h : h(xs) = x(hs) \Longrightarrow hx = xh$

- Since x is shift, A is commutative, so this is trivially verified
- Conversely, \mathcal{A} comm., x generates \mathcal{A} , then all filters are shift-inv.
- Which algebras are shift invariant (comm. & generated by single x?)
 - Infinite case: series in x or polynomials in x
 - Finite dimensional case: *polynomial algebras*, *p*(*x*) polyn. deg *n*

 $\mathcal{A} = \mathbb{C}[x]/p(x) = \{ \text{polyn. deg } < n, \text{ addition, mult. mod } p(x) \}$

Signal Model: finite dimensional case

$$\begin{array}{l} (\mathcal{A} = \mathbb{C}[x]/p(x), \mathcal{M} = \mathbb{C}[x]/p(x), \Phi) \\ \Phi : V = \mathbb{C}^n \to \mathcal{M} \end{array}$$



Algebraic Theory: Infinite Time

Realization of signal model (infinite time):

• Time marks and shift operator (Kalman 68): $q \diamond t_n = t_{n+1}$

• *k*-fold shift: $q_k \diamond t_n = t_{n+k}, \Longrightarrow q_k = q^k$

- Linear extension:
 - Extend q from t_n to set $\{\sum s_n t_n\}$
 - Extend from q^k to set of all formal sums $\left\{ \sum h_k q^k \right\}$
 - Realization: set q = x and $\diamond = \cdot \Longrightarrow t_{n+1} = x \cdot t_n$
 - Two-term recursion solution: $t_0 = 1, t_n = x^n$

$$\Longrightarrow \mathcal{M} = \{s = \sum s_n x^n\}$$
 and $\mathcal{A} = \left\{h = \sum h_k x^k\right\}$

$$\Phi:\mathbf{s}\mapsto \sum_{n\in\mathbb{Z}}x^n$$

Remark: we use *x* rather than z^{-1}



Algebraic Theory: Finite Time

Realization of signal model (finite time):

• Problem: $x \cdot x^{n-1} = x^n \notin \mathbb{C}_n[x] = \left\{ \sum_{i=0}^{n-1} s_i x^i \right\}$ $x^{-1} \cdot x^0 = x^{-1} \notin \mathbb{C}_n[x] = \left\{ \sum_{i=0}^{n-1} s_i x^i \right\}$

Boundary condition and signal extension:

•
$$x^n = r(x) = \sum_{0 \le k < n} \beta_k x^k$$
, or $x^n - r(x) = 0$

Equivalent to right b.c. $s_n = \sum_{0 \le k < n} \beta_k s_k$

Replaces vector space $\mathbb{C}_n[x]$ by $\mathcal{M} = \mathbb{C}[x]/(x^n - r(x))$

$$x^{-1} \in \mathcal{A} \text{ iff } x \not\mid r(x) \Longrightarrow x^{-1} = -\frac{1}{\beta_0} \left(\beta_1 + \beta_2 x + \dots + \beta_{n-2} - x^{n-1} \right)$$

b.c. Right and left signal extension

Signal model:

$$\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/(x^n - r(x)), \Phi(\mathbf{s}) = \sum_{0 \le k < n} s_k x^k \in \mathcal{M}$$

Electrical & Computer

Monomial signal extension: $r(x) = a \Longrightarrow x^n = a$

Finite Time and DFT

■ Signal model: $(\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/(x^n - 1), \sum_{0 \le k < n} s_k x^k \in \mathcal{M})$ ■ Fourier transform: DFT

$$\Delta : \mathcal{M} = \mathbb{C}[x]/(x^n - 1) \to \bigoplus_{0 \le k < n} \mathbb{C}[x]/(x - \omega_n^k)$$
$$\omega_n = e^{-2\pi j/n}$$

• In matrix format: $\mathsf{DFT}_n = [\omega_n^{k\ell}]_{0 \leq k, \ell < n}$



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Space Signal Model: Space Shift

Shift: symmetric definition $q \diamond t_n = (t_{n+1} + t_{n-1})/2$ *k*-fold shift: $q_k \diamond t_n = (t_{n+k} + t_{n-k})/2, \Longrightarrow q_k = q^k$

Differences wrt time model: $q_k \neq q^k$, $q_k = q_{-k}$

Lemma: The *k*-fold space shift operator is $q_k = T_k(q)$, the Chebyshev polynomials of the 1st kind

Linear extension: extend operation of q to

$$\mathcal{A} = \left\{ \sum_{k \ge 0} h_k q_k = \sum h_k T_k(q) \right\}, \quad \mathcal{M} = \left\{ \sum s_n t_n \right\}$$

Realization: q = x and $\diamond = \cdot$

time marks t_n , i.e., satisfy $x \cdot C_n = (C_{n+1} + C_{n-1})/2$ $\implies C_n$ sequence of Chebyshev polyn., $\{T_n, U_n, V_n, W_n\}$



Signal Model: Infinite Space

Signal Model:

$$\left(\mathcal{A} = \left\{h = \sum_{k \ge 0} h_k T_k\right\}, \mathcal{M} = \left\{s = \sum_{n \ge 0} s_n C_n\right\}, \Phi\right)$$

C-transform: $\Phi : V = \ell^2(\mathbb{N}) \to \mathcal{M} = \left\{\sum_{n \ge 0} s_n C_n\right\}$
$$\Phi(\mathbf{s}) = \sum_{n \ge 0} s_n C_n, \ C_n \in \{T_n, U_n, V_n, W_n\}$$

Follows from property of Chebyshev polyn.: *k*-fold shift

$$T_k \cdot C_n = \left(C_{n+k} + C_{n-k}\right)/2$$

Fourier transform: DSFT, e.g., choose $C_n = T_n$

$$\Delta: \qquad \mathcal{M} \rightarrow \left(\bigoplus_{\omega \in [0,\pi]} \mathbb{C}\right) = \mathbb{C}^{[0,\pi]}$$
$$s = S(x) \mapsto S(\cos \omega)_{\omega \in [0,\pi]} = \omega \mapsto S(\cos \omega).$$



Signal Model: Finite Space

• Left b.c.: $C_1 = ax + b, a \neq 0, \Longrightarrow C_{-1} = \frac{2-a}{a}C_1 - \frac{2b}{a}C_0$

- Monomial signal extension: $C \in \{T, U, V, W\}$
- Right b.c.: problem with $T_1 \cdot C_{n-1} = x \cdot C_{n-1} = (C_{n-2} + C_n)/2 \notin \left\{ \sum_{0 \ge k < n} s_k C_k \right\}$
- Lemma (Monomial right sig. extension): Let $C \in \{T, U, V, W\}$ Only 4 right bc yield monomial right sig. ext. for $\mathcal{M} = \mathbb{C}[x]/p(x)$ $C_n = C_{n-2}, C_n = 0, C_n = \pm C_{n-1}$ $\implies p \in \{C_n - C_{n-2}, C_n, C_n \pm C_{n-1}\}$ Electrical & Computer \implies 16 possibilities

Finite Sp.Signal Model: Finite C-transf. & DTTs Let seq. Chebyshev poly.: $C_0, \dots, C_{n-1}, C \in \{T, U, V, W\}$

Let:
$$p \in \{C_n - C_{n-2}, C_n, C_n \pm C_{n-1}\}$$

	$C_n - C_{n-2}$	C_n	$C_n - C_{n-1}$	$C_n + C_{n-1}$
T_n	$2(x^2-1)U_{n-2}$	T_n	$(x-1)W_{n-1}$	$(x+1)V_{n-1}$
	$2T_n$			
V_n	$2(x-1)W_{n-1}$	V_n	$2(x-1)U_{n-1}$	$2T_n$
	$2(x+1)V_{n-1}$			$2(x+1)U_{n-1}$

• 16 finite space signal models: $(\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/p(x), \Phi)$

Finite C-transform: $\Phi: V = \mathbb{C}^n \to \mathcal{M} = \mathbb{C}[x]/p(x)$ Electrical & Computer $\Phi(s) = \sum_{0 \le k < n} s_k C_k$

Finite Sp. Sig. Model: Finite C-transf. & DTTs

Fourier transforms: 16 DTTs (8 DCTs and 8 DSTs)

Example: Signal model for DCT, type 2

• Left bc: $s_{-1} = s_0 \Longrightarrow C_1 = C_0 = 1$ afforded by C = V

• Right bc:
$$s_n = s_{n-1} \Longrightarrow C_n = C_{n-1}$$

 $\implies p = C_n - C_{n-1} = V_n - V_{n-1} = 2(x-1)U_{n-1}$ Sig. model for DCT, type 2:

$$(\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/(x-1)U_{n-1}, \Phi)$$

DCT, type 2:

 $\Delta: \mathbb{C}[x]/(x-1)U_{n-1} \rightarrow \bigoplus_{0 \le k < n} \mathbb{C}[x]/(x-\alpha_k)$ Zeros of $(x-1)U_{n-1}, \alpha_k = \cos k\pi/n, \ 0 \le k < n$ $\mathsf{DCT-2}_n = \mathsf{diag}_{0 \le k < n} (\cos k\pi/(2n)) \cdot \left[\frac{1}{\cos k\pi/(2n)} \cdot \cos k(\ell+1/2)\pi/n\right]$ Electrical of Computer FINGINFERING

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Fast Algorithms: DTTs

DTTs: DCT, type 2: Direct sum: fast alg. Via poly. factorization $\mathbb{C}[x]/(x-1)U_{n-1}$ with basis $b = (V_0, \dots, V_{n-1})$ Property of U: $U_{2m-1} = 2 \cdot U_{m-1} \cdot T_m$ $\mathbb{C}[x]/(x-1)U_{2m-1} \cong \mathbb{C}[x]/(x-1)U_{m-1} \oplus \mathbb{C}[x]/T_m$ $DCT-2_{2m} = L_m^{2m} \cdot (DCT-2_m \oplus DCT-4_m) \cdot B.$ $B = \begin{bmatrix} 1 & & & 1 \\ & \cdot & & \cdot & \cdot \\ & & 1 & 1 & & \\ 1 & & & & -1 \\ & \cdot & & \cdot & \cdot \\ & & 1 & -1 & & \end{bmatrix} = \begin{bmatrix} I_m & J_m \\ I_m & -J_m \end{bmatrix} = (DFT_2 \otimes I_m)(I_m \oplus J_m)$



Fast Algorithms: DTTs

THEOREM 8.6. The following recursive algorithms for DTTs are based on the rational factorization $U_{2m-1} = 2 \cdot U_{m-1} \cdot T_m$. We also indicate where they first appeared in the literature (to our best knowledge).

(i) DCT-
$$1_{2m+1} = P_{2m+1} \cdot (DCT-1_{m+1} \oplus DCT-3_m) \cdot B_{2m+1}$$
, [27].
(ii) DST- $1_{2m-1} = P_{2m-1} \cdot (DST-3_m \oplus DST-1_{m-1}) \cdot B_{2m-1}$, [55].
(iii) DCT- $2_{2m} = P_{2m} \cdot (DCT-2_m \oplus DCT-4_m) \cdot B_{2m}$, [7].
(iv) DST- $2_{2m} = P_{2m} \cdot (DST-4_m \oplus DST-2_m) \cdot B_{2m}$, [52].

THEOREM 8.7. The following recursive algorithms for DTTs are based on the rational factorization $U_{2m} = V_m W_m$.

(i) DCT-
$$1_{2m} = P_{2m} \cdot (DCT-5_m \oplus DCT-7_m) \cdot B_{2m}$$
.
(ii) DST- $1_{2m} = P_{2m} \cdot (DST-7_m \oplus DST-5_m) \cdot B_{2m}$.
(iii) DCT- $2_{2m+1} = P_{2m+1} \cdot (DCT-6_{m+1} \oplus DCT-8_m) \cdot B_{2m+1}$.
(iv) DST- $2_{2m+1} = P_{2m+1} \cdot (DST-8_{m+1} \oplus DST-6_m) \cdot B_{2m+1}$.



Finite Signal Models in Two Dimensions

Visualization (without b.c.)	Signal Model $\mathcal{A}=\mathcal{M}$	Fourier Transform
time, separable	$\mathbb{C}[x,y]/\langle x^n{-}1,y^n{-}1 angle$ time shifts: x, y	$\mathrm{DFT}_n\otimes\mathrm{DFT}_n$
space, separable	for example $\mathbb{C}[x,y]/\langle T_n(x),T_n(y) angle$ space shifts: x, y	$ ext{DCT}_n \otimes ext{DCT}_n$ (16 types)



time, nonseparable	$\mathbb{C}[u,v]/\langle u^n-1,u^{rac{n}{2}}-v^{rac{n}{2}} angle$ time shifts: u, v	$ ext{DDQT}_{n imesrac{n}{2}}$ Püschel ICASSP '05 (separable, Mersereau)
space, nonseparable	$\mathbb{C}[u,v,w]/\langle T_{n/2}(u),T_{n/2}(v),\ 4w^2-(u+1)(v+1) angle$ space shifts: u, v, w	$\mathrm{DQT}_{n imes rac{n}{2}}$ Püschel ICIP '05
space, nonseparable Electrical & Computer ENGINEERING	$\mathbb{C}[x,y]/\langle C_n(x,y),\overline{C}_n(x,y) angle$ space shifts: u, v	DTT $_{n imes n}$ Püschel ICASSP '04

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